

Key:

Name:

Student number:

Computational Science 260

Second Midterm Exam

Nov. 30, 1995

Marks

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1. Write the predicate $\text{max}(L, M)$ in Prolog. Here, L is a list of numbers, and $\text{max}(L, M)$ must succeed if M is the largest element in the list. Otherwise, the predicate should fail. The predicate should also fail for the empty list.

$\text{max}([X], X).$

$\text{max}([X | \text{Tail}], X) :- \text{max}(\text{Tail}, Z), X > Z.$

$\text{max}([X | \text{Tail}], Z) :- \text{max}(\text{Tail}, Z), X \leq Z.$

- CHS 2. Let $A = \mathcal{P}\{3\}$. Give A in roster notation, and find $\#A$.

$A = \{ \{ \}, \{ 3 \} \}$

$\#A = 2$

\mathcal{P} Powerset

$\#A$ Cardinality

$A = \{ \emptyset, \{ 3 \} \}$

$\#A = 2$

$A = \{ \{ \}, \{ 3 \} \}$

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- CH 3. Let $f : X \rightarrow Y$ be a partial function from X to Y . Use appropriate phrases to characterize f under the following conditions

(a) $\text{dom } f = X$

: total function

(b) $\text{dom } f = X, \text{ran } f = Y$

: surjective

(c) $\text{dom } f = X, \text{ran } f = Y, x \neq y \Rightarrow f(x) \neq f(y)$

: bijective

injective

ONTO

///

Domain is whole
Range is a subset
of Range

- CHS 4. Let A be a set of people who have attended party 1, and let B be the set of people who have attended party 2. Furthermore, C be the set of people that have attended both parties, and let D be the set that have one, but not both. Use the normal set operations, such as union, intersection, etc, to express C and D in terms of A and B .

Tricky $C = A \cap B$

Tricky $D = (A \cup B) - (A \cap B)$

5. Two relations R and S are given as follows

CHS $R = \{(mary, john), (jane, brent), (lia, paul), (anne, ken)\}$

$S = \{(lia, carl)\}$

Find the set $A = \{(x, y) \mid x \notin \text{dom } S \wedge xRy\} \cup S$ in roster notation.

$A = \{(mary, john), (jane, brent), (lia, carl), (anne, ken)\}$

A updates all information about lia.

- CHS 6. Let S be the sibling relation, that is, $(x, y) \in S$ iff x and y have both parents in common. Let H be the halfsibling relation, that is $(x, y) \in H$ iff x and y share the father or the mother, but not both. Furthermore, let I be the identity relation.

- (a) Prove that $S \cup I$ is an equivalence relation by verifying that all the properties required for an equivalence relation are met.
(b) Is $H \cup I$ an equivalence relation? Check all the properties required for a relation to be an equivalence relation, and indicate which ones are met.

- a) 1. $S \cup I$ is reflexive because I is
2. $S \cup I$ is symmetric: If x is sibling of y ,
 y is sibling of x
3. $S \cup I$ is transitive.

- b) 1. $H \cup I$ is reflexive because I is
2. $H \cup I$ is symmetric

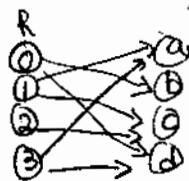
3. $H \cup I$ is not transitive: A halfsibling of a halfsibling is normally not a half-sibling.

- ① Find Relation
- ② D/T Matrix
- ③ Calcul

CH 6

7. Let $R: 0..3 \rightarrow \{a, b, c, d\}$ be a relation, and let the relation matrix of R be given as follows

$$M_R = \begin{matrix} & \begin{matrix} a & b & c & d \end{matrix} \\ \begin{matrix} 0 \\ 1 \\ 2 \\ 3 \end{matrix} & \begin{bmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 1 \end{bmatrix} \end{matrix}$$

 $\sim R$

Look backwards

$\{(a,1), (a,3), (b,0), (c,1), (c,2), (d,0), (d,3)\}$

$R \circ R \sim \{(a,b), (b,a), (a,c), (c,a), (d,c), (c,d)\}$

Give the relation $R \circ R \sim$ in roster notation.

$$M_R = \begin{matrix} & \begin{matrix} a & b & c & d \end{matrix} \\ \begin{matrix} 0 \\ 1 \\ 2 \\ 3 \end{matrix} & \begin{bmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 1 \end{bmatrix} \end{matrix}$$

 M_R

$$M_{R \sim} = \begin{matrix} & \begin{matrix} 0 & 1 & 2 & 3 \end{matrix} \\ \begin{matrix} a \\ b \\ c \\ d \end{matrix} & \begin{bmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 1 \end{bmatrix} \end{matrix}$$

 $M_{R \sim}$

in roster notation:

$\{(0,0), (0,3), (1,1), (1,2), (1,3), (2,1), (2,2), (3,0), (3,1), (3,3)\}$

8. Let $real$ be a basic type of Z , and let $x: real, y: Z$. Give the Z declaration of the function f defined by $f(x, y) = x + y$. Here, $x + y$ is evaluated as in Pascal, that is, mixed mode expressions yield a real result.

$f: real \times Z \rightarrow real$

9. Consider the following Z fragment which implements a phone directory.

$[name, phone]$

$message ::= ok | not_in_directory$

$book$
$directory : name \rightarrow phone$
$subscribers = \text{dom } directory$

subscriber

- (a) Write two schemas for finding the phone number of a subscriber. The first of these two schemas should apply for the case where the name, call it $x?$, is in the directory, and the second schema should deal with the case where $x?$ cannot be found in the directory.
- (b) Suppose the declaration of *directory* is changed from $directory : name \rightarrow phone$ to $directory : name \rightarrow \mathcal{P} phone$. Write a schema for this case, in which all phone numbers of $x?$ are output.

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a)

find_number	not_listed
$\exists book$ $number! : phone$ $x? : name$ $confirmation! : message$	$\exists book$ $x? : name$ $confirmation! : message$
$x? \in subscribers$ $number! = directory\ x$ $confirmation! = ok$	$x? \notin subscribers$ $confirmation! = not_in_directory$

b)

find_numbers
$\exists book$ $numbers! : \mathcal{P} phone$ $x? : name$ $confirmation! : message$
$numbers! = \{y : phone \mid x? \mapsto y \in directory\}$

Also possible:
 $numbers! = directory \cap \{x?\} D.$